

Remotely Controlling TrustZone Applications? A Study on Securely and Resiliently Receiving Remote Commands

Shengye Wan, Kun Sun, Ning Zhang, Yue Li



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OUTLINE

- Introduction
 - System Overview
 - System Evaluation
 - Takeaways
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INTRODUCTION

Background: Mobile Device Management

- Mobile device management (MDM)
 - Enable corporate administrators to remotely perform essential functions
 - Supportability, security, and corporate functionality

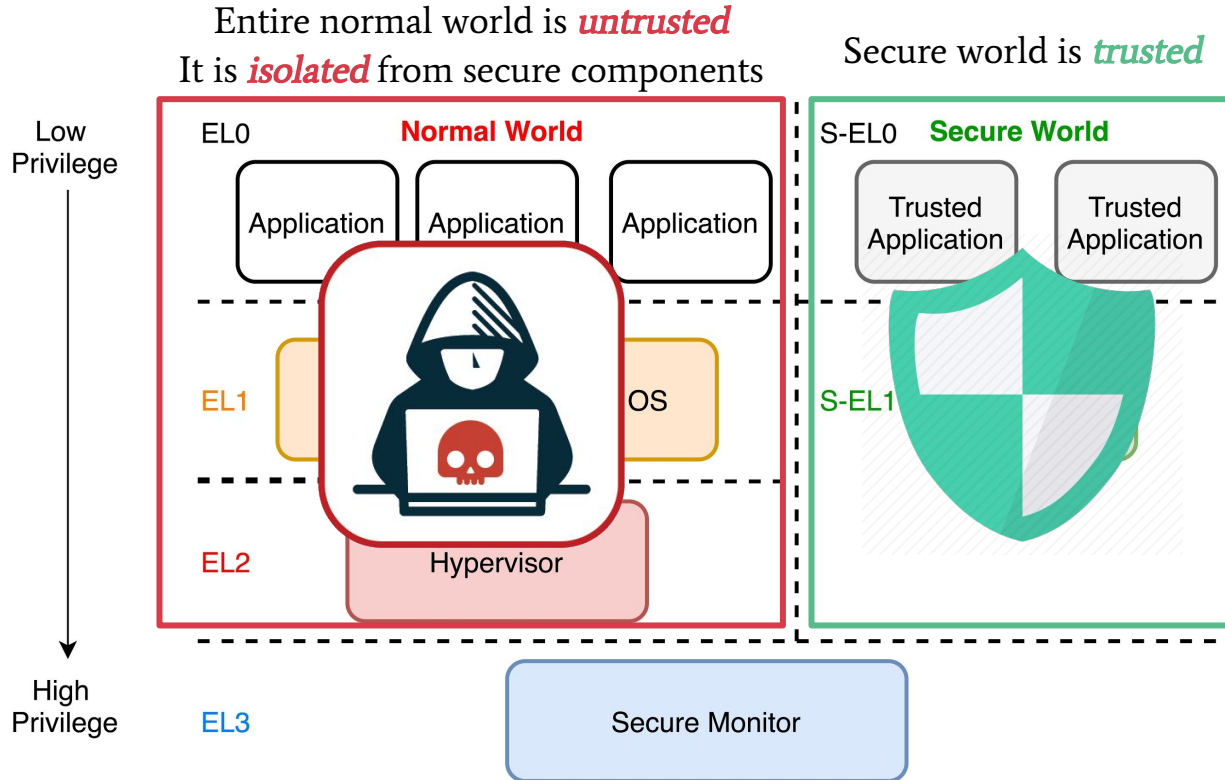


Security of MDM Agents

- MDM workflow
 - Administrator <-> Management Commands <-> MDM Agents (clients)
- MDM agents are security-sensitive
 - Rich OS cannot be trusted to hold MDM agents
 - 859 CVEs are reported in 2020 for Android [1]
 - Opportunities to enhance MDM agents' security



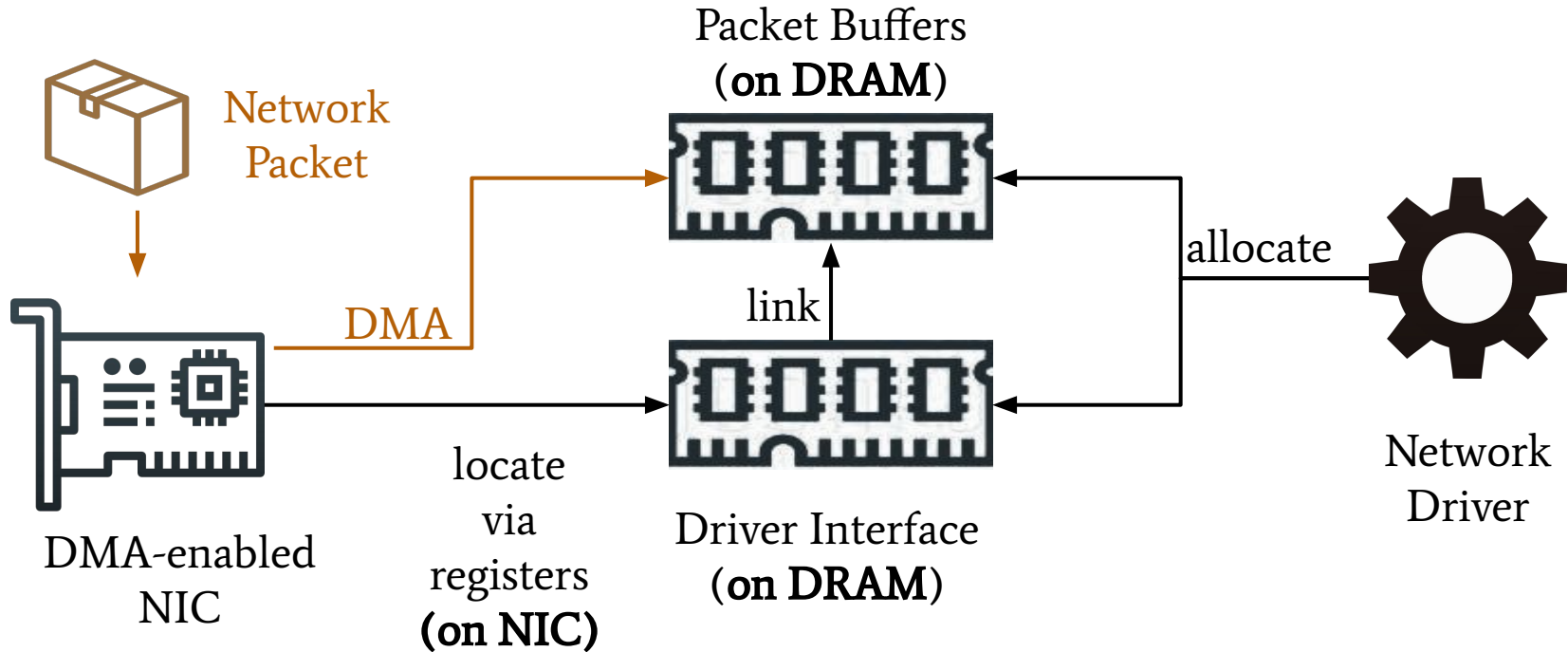
Background: ARM TrustZone Technology



Motivation: Two Worlds Need to Share One NIC

- MDM agents require network service
 - Remote attestation, remote control, remote troubleshooting
- Secure world (SW) does not have an exclusive NIC
 - Commercial devices only equip one set of network devices
 - Limited hardware spaces on mobile
 - NW and SW need to share the NIC
- **Question:** With a shared Network Interface Card (NIC), how to provide a reliable network for ARM TrustZone secure world?

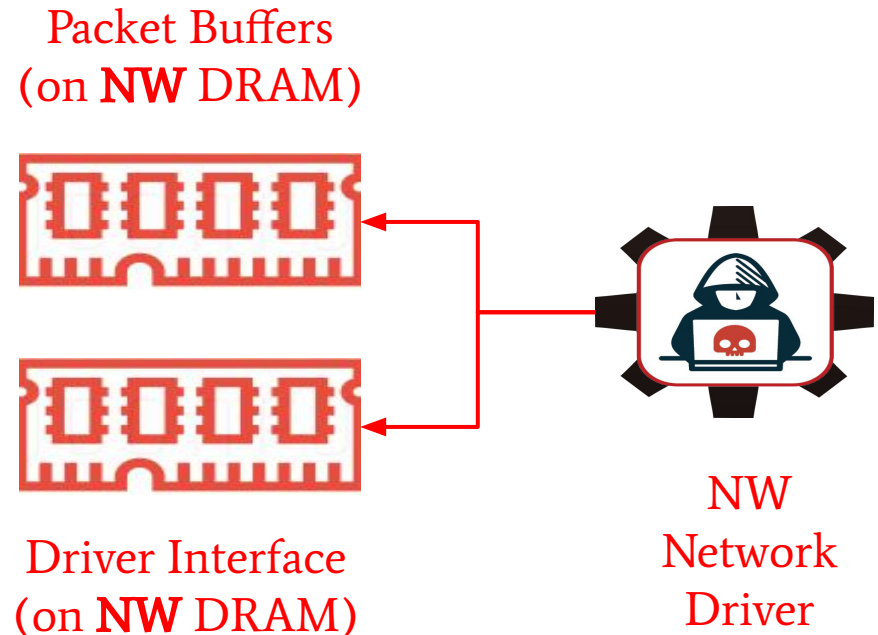
Background: NIC Workflow



SYSTEM OVERVIEW

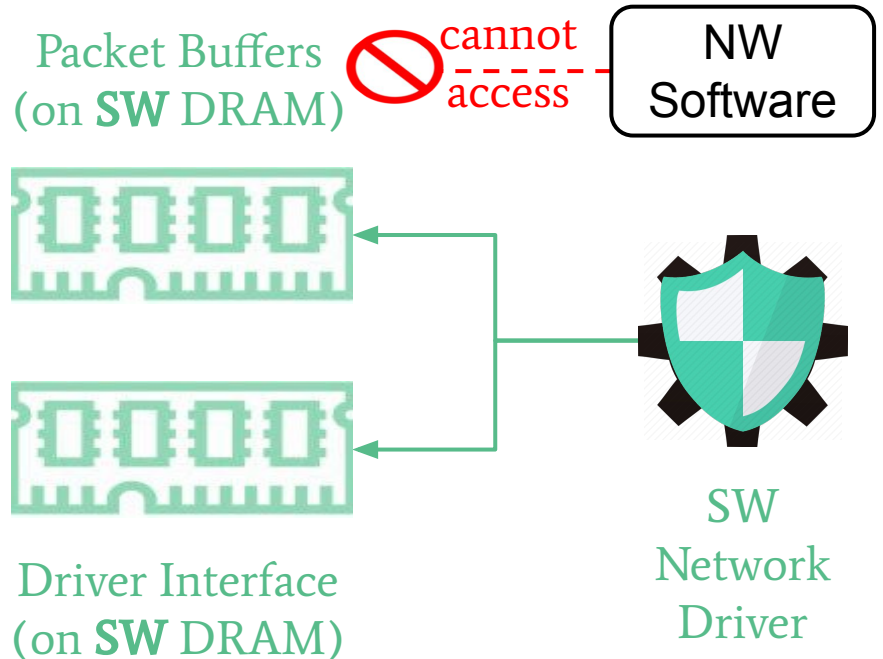
How to Share One NIC Between Two Worlds?

- Option-1: sharing the single network driver in NW
 - Pros: providing good normal world performance
 - Cons: not reliable for the secure world



Sharing One NIC: Option-2

- Option-2: sharing the single network driver in SW
 - Pros: reliable for SW
 - Cons: introducing large overhead
 - NW software cannot access packet buffers directly



Sharing One NIC: Option-3

- Option-3: deploying two network drivers in each world

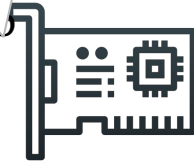
- Pros: reliable and performance

- Cons: very difficult to schedule two drivers

- One NIC only connects to one driver's interface

None of these options works!

NW Driver Interface + Packet Buffers



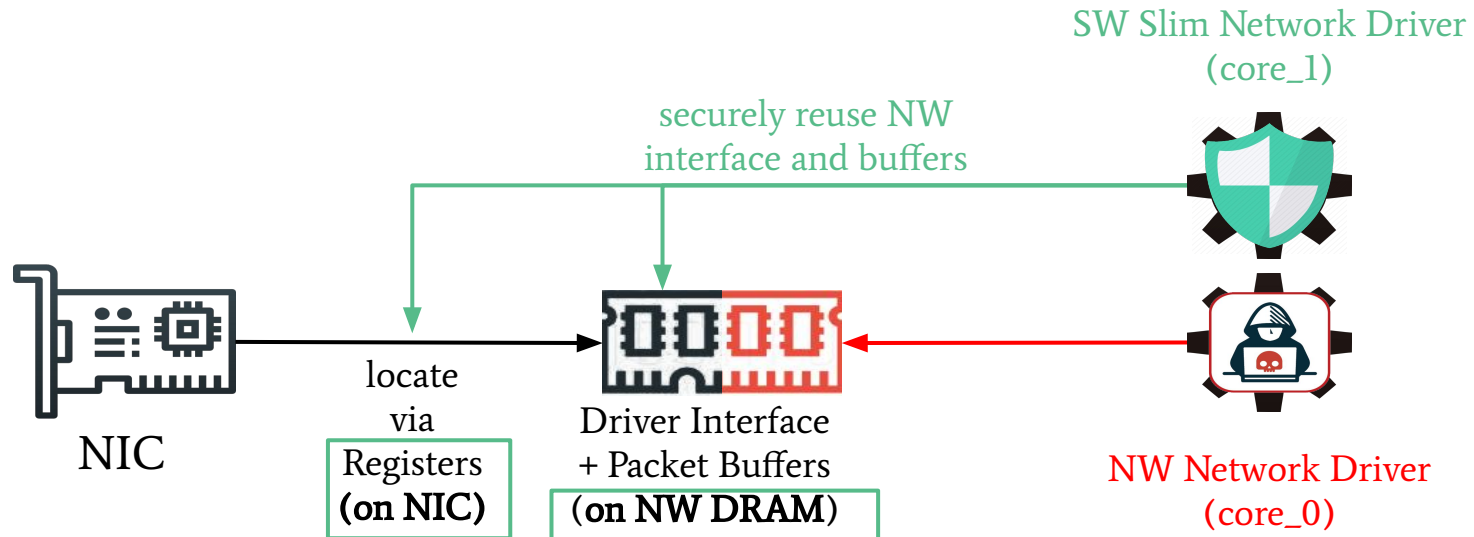
NIC



NW Driver Interface + Packet Buffers

Our Solution: TZNIC

- Deploying a complete NW-driver and a slim SW-driver
 - **Key idea:** executing two drivers simultaneously on the multi-core platform
 - Multiplexing the NW-driver's interface



TZNIC Challenges

1. Filling the semantic gap to use NW-driver's interface reliably
 - SW-driver should not put any trust in the normal world
 - SW-driver should not require any collaboration from the normal world
2. Resisting interference from the normal world
 - Securely sharing the interface and buffers with NW-driver

Resolving Challenge-1: Filling Semantic Gap

- Locating NW driver's interface via the NIC registers
 - Registers indicate the ring buffer information
 - Registers are readable to the secure world
- Locating the packets via the NW driver's interface
 - Interface and buffers are saved in the NW memory
 - Secure world has the privilege to read/write
 - NW driver uses fixed-format interface to communicate with NIC
- Does not request any collaboration of the normal world

Resolving Challenge-2: Resisting NW interference

- Reading packets in parallel of NW-driver
 - SW-driver wakes itself periodically to receive the packets
 - One receiving buffer can be read by two drivers simultaneously
- Saving the secure-world packets to the secure memory
 - Each buffer should be independent and loss-tolerant (e.g., UDP)
 - Normal-world attacker cannot access

SYSTEM EVALUATION

TZNIC Implementation

- Implementing our prototype based on ARM-TF [2]
 - Marvell Yukon-II NIC & Marvell sky-2 driver (v 1.30)

- TZNIC's slim driver's size is 18.63% of the original driver
 - Full-fledged normal-world sky-2 driver: 5707 LOC
 - TZNIC slim secure-world driver: 1063 LOC


TZNIC Evaluation - Reliability

- Attacker capacity
 - Brute-force deleting the packet from a specific IP
 - Benchmark iPerf [3] cannot receive any packet under our interference
- Under the interference of our attacker
 - TZNIC receives 67% of the packets on average
 - 22% - 92%

TAKEAWAYS

Summary


1. We can support software in TrustZone secure world with reliable network
2. Secure-world driver can reliably reuse the normal-world driver's interface
 - a. Secure world has higher privilege to inspect on-device registers
 - b. Secure world has higher privilege to read normal-world driver's data
 - c. Secure world has higher privilege to get activated
3. TZNIC makes 0 modifications or requirements on the rich OS



Thanks & Questions?

Shengye Wan

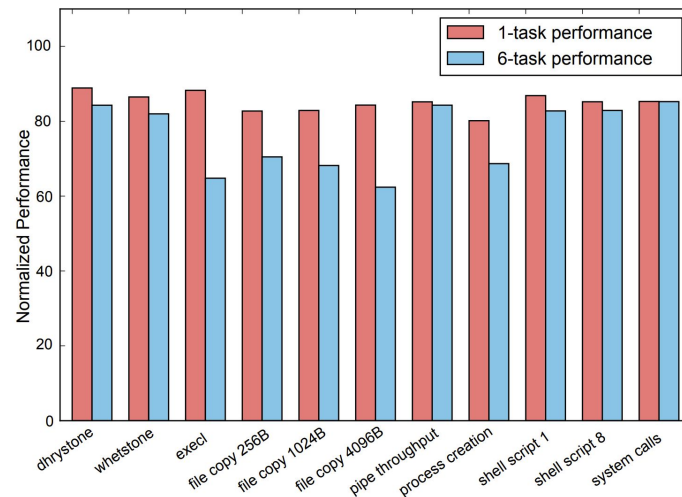
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Backup Slides

TZNIC Evaluation - Rich OS Overhead

- When TZNIC wakes up, rich OS will suffer 16.7% overhead
- The overall overhead can easily be improved
 - TZNIC does not wake up often
 - The wake-up frequency can be adjusted
 - To promise 95% of the rich OS performance:
 - TZNIC wakes 10ms among every 80ms



Future Works

1. Protecting network devices from Denial-of-Service attacks
 - Configuring the NIC as a secure-world hardware
2. Deploying multiple TZNIC in secure world
 - Solution-1: moving TZNIC into secure application layer
 - Solution-2: Using new ARM TrustZone feature
 - Achieve virtualization in the secure world

Background: Cross-World Context Switch

- SMC
 - ARM special instruction to enter the Secure Monitor (EL3) code
 - *Core-i* can only use SMC to switch the status of *core-i*
- Interrupt
 - SW-interrupt is promised to route to secure world
 - Interrupt untrusted NW execution
 - One interrupt may arrive on
 - One specific *core-i* (Private Peripheral Interrupt)
 - Multiple cores (Shared Peripheral Interrupt, Software Generated Interrupts)
 - NW-interrupt can get handled in both worlds